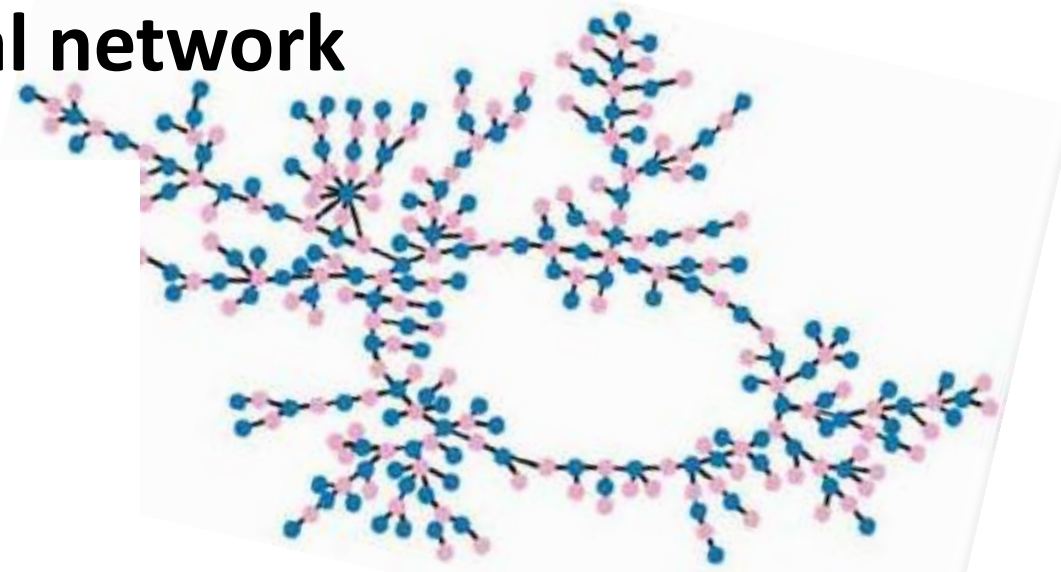
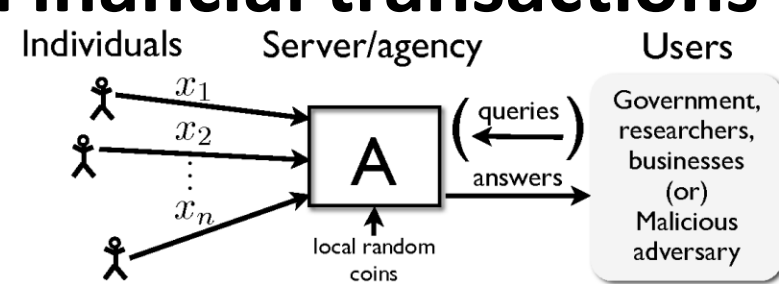


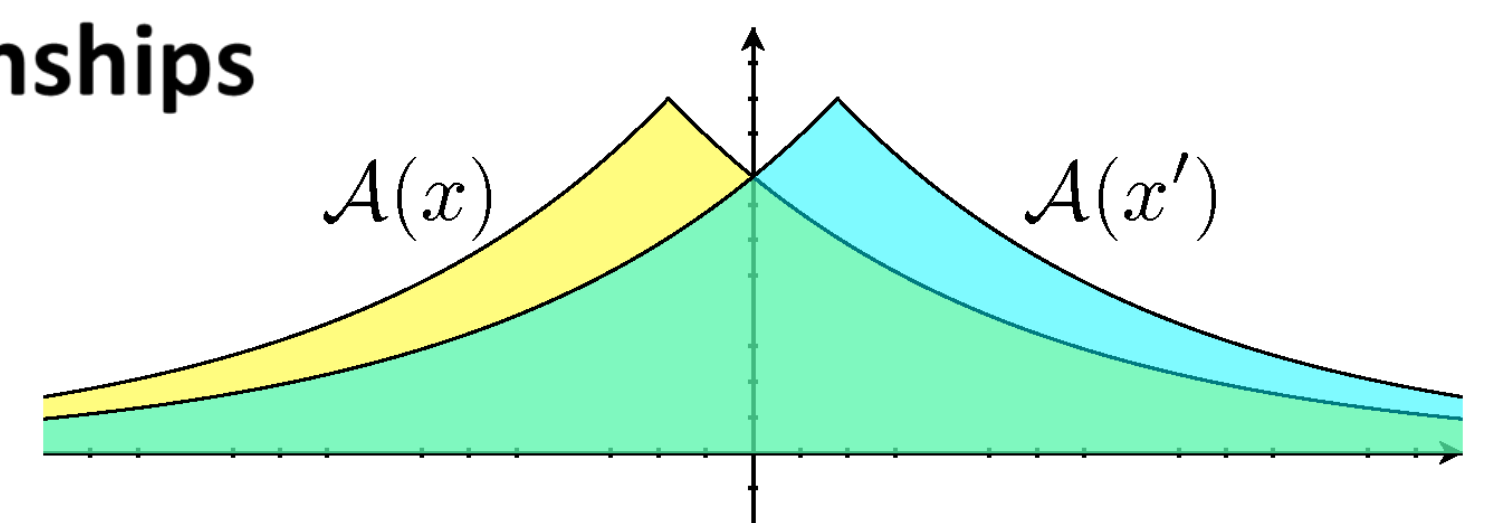
## Network Data

- Many data sets can be naturally represented as network of pairwise relationships (graph)
  - Friendship in online social network
  - Romantic relationships
  - Financial transactions
- Goal: Publish info about network structure for use by researchers, gov't agencies, ...
- Privacy concerns: Sensitive information
  - De-anonymizing social networks [BDK'07, NS'09]
  - Link prediction by de-anonymization [NSR'11]
- This work: new privacy-preserving algorithms



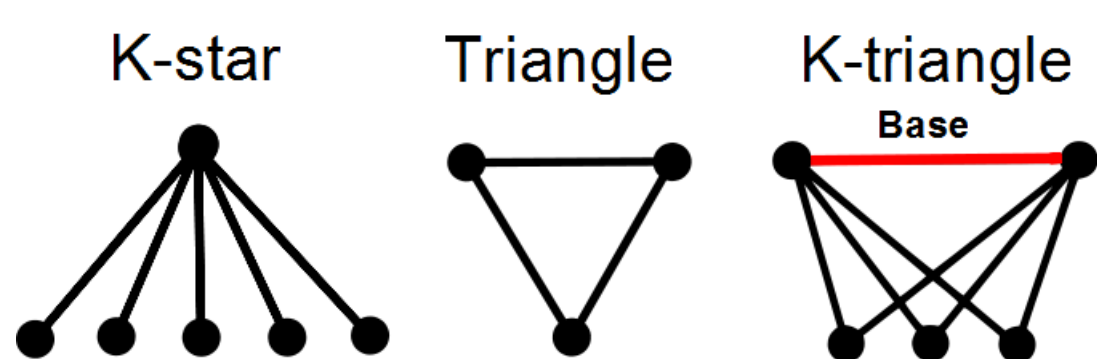
## Measuring Privacy

- Goal: Rigorous privacy guarantee
- Challenges
  - Attacker's prior information unknown
  - Attack algorithm / methodology unknown
- $\epsilon$ -Differential privacy [DMNS'06]: For all neighbors  $G$  and  $G'$  (differing in one record):
- Hides the presence/absence of individual relationships



## Subgraph Counts

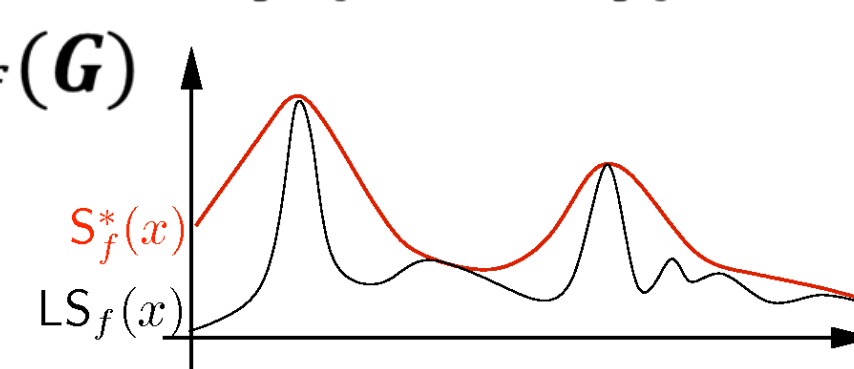
- For graphs  $G$  and  $H$ ,  $f_H(G) = \#$  of occurrences of  $H$  in  $G$



- Subgraph counts are widely used in:
  - Exponential random graph models
  - Descriptive graph statistics (e.g. clustering coefficient)
- Our work: We present d.p. algorithms for accurately releasing subgraph counts

## Techniques

- Add noise, proportional to  $S/\epsilon$  (where  $S$  is sensitivity). How to minimize  $S$ ?
- This work: Instance-specific noise
  - Local sensitivity (also used in [RHMS'09])
$$LS_f(G) = \max_{G' \in N(G)} |f(G) - f(G')|$$
- Nearly optimal amplitude, but not d.p.
- Previous work: Global sensitivity (mostly)
 
$$GS_f(G) = \max_G LS_f(G)$$
- Main approaches:
  - Smooth sensitivity [NRS'07]
  - D.p. approximation to  $LS_f(G)$  (this work).



## Results

- New algorithms:
  - K-stars (linear time algorithms for computing smooth sensitivity)
  - K-triangles (d.p. approximation to  $LS_f(G)$ ).
- Bounds on noise in terms of  $LS_f(G)$ :
  - Conditions, when amplitude of noise is close to  $\frac{LS_f}{\epsilon}$
- Average-case analysis in Gilbert model  $G(n,p)$  ( $np = \omega(\sqrt{n \log n})$ ,  $\epsilon = o(\log n)$ ):
 
$$\frac{|A_H(G) - f_H(G)|}{f_H(G)} = O\left(\frac{1}{\epsilon n^2 p}\right)$$
- Computing smooth sensitivity is NP-hard for some subgraph queries (K-triangles, K-cycles)

## Experimental Evaluation

Results of experimental evaluation for  $\epsilon = 0.5$ :

Dataset	$H$	Median Relative Error
CaGrQc	$\Delta$	0.081
Nodes = 5242	$2\star$	0.0041
Edges = 2890	$3\star$	0.144
CaHepTh	$\Delta$	0.11
Nodes = 9877	$2\star$	0.002
Edges = 51971	$3\star$	0.0009
CaAstroPh	$\Delta$	0.017
Nodes = 18772	$2\star$	0.00043
Edges = 396160	$3\star$	0.0023
CaHepPh	$\Delta$	0.007
Nodes = 12008	$2\star$	0.00038
Edges = 237010	$3\star$	0.001
CaCondMat	$\Delta$	0.077
Nodes = 23133	$2\star$	0.00158
Edges = 186936	$3\star$	0.011
Enron	$\Delta$	0.032
Nodes = 36692	$2\star$	0.00055
Edges = 267662	$3\star$	0.0021