

# CIS 700: “algorithms for Big Data”

## Lecture 11: K-means

Slides at <http://grigory.us/big-data-class.html>

**Grigory Yaroslavtsev**

<http://grigory.us>



# K-means Clustering

- Given  $X = \{x_1, \dots, x_n\} \in \mathbb{R}^d$  find a set of centers  $C = (c_1, \dots, c_k)$  that minimizes

$$\sum_{x \in X} \min_{i \in [k]} \|x - c_i\|^2$$

- NP-hard problem
- Popular heuristic local search (Lloyd's alg.)
- For a fixed partitioning  $P_1, \dots, P_k$ :

$$c_j = \frac{1}{|P_j|} \cdot \sum_{i \in P_j} x_i$$

# Dimension reduction for K-means

- Let  $cost_P(X) = \inf_c cost_{P,c}(X)$
- For  $0 < \epsilon < \frac{1}{2}$  let  $f: X \rightarrow \mathbb{R}^n$  be such that  
 $\forall i, j: (1 - \epsilon) \|x_i - x_j\|_2^2 \leq \|f(x_i) - f(x_j)\|_2^2 \leq (1 + \epsilon) \|x_i - x_j\|_2^2$
- $\hat{P}$  is a  $\gamma$ -approx. clustering for  $f(X)$
- $P^*$  is an optimal clustering for  $X$
- **Lemma.**

$$cost_{\hat{P}} \leq \gamma \left( \frac{1 + \epsilon}{1 - \epsilon} \right) cost_{P^*}(X)$$

# Dimension reduction for K-means

- Let  $cost_P(X) = \inf_c cost_{P,c}(X)$
- For  $0 < \epsilon < \frac{1}{2}$  let  $f: X \rightarrow \mathbb{R}^{d'}$  be such that  
 $\forall i, j: (1 - \epsilon) \|x_i - x_j\|_2^2 \leq \|f(x_i) - f(x_j)\|_2^2 \leq (1 + \epsilon) \|x_i - x_j\|_2^2$
- $\hat{P}$  is a  $\gamma$ -approx. clustering for  $f(X)$
- $P^*$  is an optimal clustering for  $X$
- **Lemma.**

$$cost_{\hat{P}} \leq \gamma \left( \frac{1 + \epsilon}{1 - \epsilon} \right) cost_{P^*}(X)$$

- $d' = O\left(\log \frac{n}{\epsilon^2}\right)$  suffices by the JL-lemma

# Dimension reduction for K-means

- Fix a partition  $P = (P_1, \dots, P_k)$

$$\begin{aligned}
 cost_P(X) &= \sum_{j \in [k]} \sum_{i \in P_j} \left\| x_i - \frac{1}{|P_j|} \sum_{i' \in P_j} x_{i'} \right\|_2^2 \\
 &= \sum_{j \in [k]} \frac{1}{|P_j|} \sum_{i \in P_j} \left( \sum_{i' \in P_j} \|x_i\|_2^2 - 2 \langle x_i, \sum_{i' \in P_j} x_{i'} \rangle + \left\| \sum_{i' \in P_j} x_{i'} \right\|_2^2 \right) \\
 &= \sum_{j \in [k]} \frac{1}{|P_j|} \sum_{i \in P_j} \sum_{i' \in P_j} \left( \frac{\|x_i\|_2^2 + \|x_{i'}\|_2^2}{2} - \langle x_i, x_{i'} \rangle \right) \\
 &\quad \sum_{j \in [k]} \frac{1}{2|P_j|} \sum_{i \in P_j} \sum_{i' \in P_j} (\|x_i - x_{i'}\|_2^2)
 \end{aligned}$$

- $(1 - \epsilon) cost_P(X) \leq cost_P(f(X)) \leq (1 + \epsilon) cost_P(X)$
- $(1 - \epsilon) cost_{\hat{P}}(X) \leq cost_{\hat{P}}(f(X)) \leq \gamma cost_{P^*}(f(X)) \leq \gamma cost_{P^*}(X)$

# K-means++ Algorithm

- First center uniformly at random from  $X$
- For a set of centers  $C$  let:

$$d^2(x, C) = \min_{c \in C} \|x - c\|_2^2$$

- Fix current set of centers  $C$
- Subsequent centers: each  $x_i$  with prob.

$$\frac{d^2(x_i, C)}{\sum_{x_j \in X} d^2(x_j, C)}$$

- Gives  $O(\log k)$ -approx. to OPT in expectation

# K-means|| Algorithm

- First center  $C$ : sample a point uniformly
- Initial cost  $\psi = \sum_x d^2(x, C)$
- For  $O(\log \psi)$  times do:
  - Repeat  $\ell$  times (in parallel)
    - $C' =$  sample each  $x_i \in X$  indep. with prob.

$$p_x = \frac{d^2(x_i, C)}{\sum_{x_j \in X} d^2(x_j, C)}$$

- $C \leftarrow C \cup C'$
- For  $x \in C$ :
  - $w_x =$  the #points belonging to this center
- Cluster the weighted points in  $C$  into  $k$  clusters

# K-means|| Algorithm

- Oversampling factor  $\ell = \Theta(k)$
- #points in  $C$ :  $\ell \log \psi$
- **Thm.** If  $\alpha$ -approx. used in the last step then  $k$ -means|| obtains an  $O(\alpha)$ -approx. to  $k$ -means
- If  $\Psi$  and  $\Psi'$  are the costs of clustering before and after one outer loop iteration then:

$$E[\Psi'] = O(OPT) + \frac{k}{e\ell} \Psi$$



# K-means|| Analysis

- For a set of points  $A = \{a_1, \dots, a_t\}$  centroid  $c_A$ :

$$c_A = \frac{1}{|T|} \sum a_t$$

- Order  $a_1, \dots, a_T$  in the increasing order by distance from  $c_A$
- Fix a cluster  $A$  in OPT
- Fix  $C$  prior to the iteration and let:

$$\phi(C) = \sum_x d^2(x, C)$$

$$\phi_A(C) = \sum_a d^2(a, C)$$

- Let  $p_t = \frac{d^2(a_t, C)}{\phi(C)}$  be the probability of selecting  $a_t$
- Probability that  $a_t$  is the smallest one chosen:

$$q_t = p_t \prod_{j=1}^{t-1} (1 - p_j)$$

# K-means|| Analysis

- Can either assign all points to some selected  $a_t$  or keep the original clustering:

$$s_t = \min \left( \phi_A, \sum_{a \in A} \|a - a_t\|^2 \right)$$

- $E[\phi_A(C \cup C')] \leq \sum_t q_t s_t + q_{T+1} \phi_A(C)$

where  $q_{T+1}$  = prob. that no point in  $A$  is selected

- Simplifying assumption: consider the case when all  $p_t = p$  (mean field analysis)
- $q_t = p(1 - p)^t$  (decreasing sequence)

# K-means|| Analysis

- $s'_t = \sum_{a \in A} \|a - a_t\|^2$
- $\{s'_t\}$  is an increasing sequence

$$\begin{aligned} \sum_t q_t s_t &\leq \sum_t q_t s'_t \\ &\leq \frac{1}{T} \left( \sum_t q_t \sum_t s'_t \right) \\ &= \left( \sum_t q_t \cdot \frac{1}{T} \sum_t s'_t \right) \\ &= \left( \sum_t q_t \cdot 2 \phi_A^* \right) \end{aligned}$$

- $E[\phi_A(C \cup C')] \leq (1 - q_{T+1}) 2 \phi_A^* + q_{T+1} \phi_A(C)$

