

Homework 4: December 04

Name: YOUR NAME HERE

Due: Monday, December 11, 11:59pm EST

Problem 4.1 (Sparse recovery) In this problem all vectors are in \mathbb{R}^n . Recall that the sparse recovery error is defined as:

$$\text{Err}^k(f) = \min_{g: \|g\|_0=k} \|f - g\|_1.$$

Give a formal argument that shows that $\text{Err}^k(f) = \sum_{i \notin S} |f_i|$ where S is the set of indices of k largest (by absolute value) entries of f .

Problem 4.2 (Dyadic intervals) Let n be a power of two. Consider the following family of partitions of the interval $1, \dots, n$ into intervals:

$$\begin{aligned} I_0 &= \{\{1\}, \{2\}, \dots, \{n\}\} \\ I_1 &= \{\{1, 2\}, \{3, 4\}, \{5, 6\}, \dots, \{n-1, n\}\} \\ I_2 &= \{\{1, 2, 3, 4\}, \{5, 6, 7, 8\}, \dots, \{n-3, n-2, n-1, n\}\} \\ &\dots \\ I_{\log n} &= \{\{1, \dots, n\}\}, \end{aligned}$$

where the partition I_k consists of intervals of length 2^k . Show that any subinterval i, \dots, j where $i \leq j$ can be represented as a disjoint union of at most $2 \log n$ intervals from the above family.

Problem 4.3 (Generating uniform distribution) Given a stream of numbers a_1, \dots, a_n where each number is an integer between 1 and m design an algorithm that scans the stream and at every point during the scan maintains a uniformly at random chosen sample of k numbers from the stream. Your algorithm should use space $O(k \log m)$.

Problem 4.4 (Bipartiteness via connectivity) Consider the following reduction: given a connected undirected graph $G(V, E)$ construct a new graph $G'(V_1 \cup V_2, E')$ where V_1 and V_2 are copies of V and for each edge $(u, v) \in E$ we create two edges (u_1, v_2) and (u_2, v_1) where u_i and v_i are copies of u and v in V_i . Prove the following two statements:

1. If G is bipartite then the number of connected components in G' equals 2.
2. If G is non-bipartite then G' is connected.